

## International Journal of Advanced Research in Science, Engineering and Technology

Vol. 6, Special Issue, August 2019

International Conference on Recent Advances in Science, Engineering, Technology and Management at Sree Vahini Institute of Science and Technology-Tiruvuru, Krishna Dist, A.P

# Secure Data Sharing In Cloud Computing Using Revocable-Storage Identity-Based Encryption

A.Muthareddy, G.V.Ramana

Assistant Professor, CSE Department Sree Vahini Institute of Science and Technology Associate Professor, CSE department, Sree Vahini Institute of Science and Technology

ABSTRACT: Cloud computing provides a flexible and convenient way for data sharing, which brings various benefits for both the society and individuals. But there exists a natural resistance for users to directly outsource the shared data to the cloud server since the data often contain valuable information. Thus, it is necessary to place cryptographically enhanced access control on the shared data. Identity-based encryption is a promising cryptographically primitive to build a practical data sharing system. However, access control is not static. That is, when some user's authorization is expired, there should be a mechanism that can remove him/her from the system. Consequently, the revoked user cannot access both the previously and subsequently shared data. To this end, we propose a notion called revocable-storage identity-based encryption (RS-IBE), which can provide the forward/backward security of ciphertext by introducing the functionalities of user revocation and ciphertext update simultaneously. Furthermore, we present a concrete construction of RS-IBE, and prove its security in the defined security model. The performance comparisons indicate that the proposed RS-IBE scheme has advantages in terms of functionality and efficiency, and thus is feasible for a practical and cost-effective data-sharing system. Finally, we provide implementation results of the proposed scheme to demonstrate its practicability. Index Terms—Cloud computing, data sharing, revocation, Identity-based encryption, ciphertext update, decryption key exposure.

#### I. INTRODUCTION

CLOUD computing is a paradigm that provides massive computation capacity and huge memory space at a low cost [1]. It enables users to get intended services irrespective of time and location across multiple platforms (e.g., mobile devices, personal computers), and thus brings great convenience to cloud users. Among numerous services provided by cloud computing, cloud storage service, such as Apple's iCloud [2], Microsoft's Azure [3] and Amazon's S3 [4], can offer a more flexible and easy way to share data over the Internet, which provides various benefits for our society [5], [6]. However, it also suffers from several security threats, which are the primary concerns of cloud users [7]. Firstly, outsourcing data to cloud server implies that data is out control of users. This may cause users' hesitation since the outsourced data usually contain valuable and sensitive information. Secondly, data sharing is often implemented in an open and hostile environment, and cloud server would become a target of attacks. Even worse, cloud server itself may reveal users' data for illegal profit. Thirdly, data sharing is not static. That is, when a user's authorization gets expired, he/she should no longer possess the privilege of accessing the previously and subsequently shared data. Therefore, while outsourcing data to cloud server, users also want to control access to these data such that only those currently authorized users can share the outsourced data. A natural solution to conquer the aforementioned problem is to use cryptographically enforced access control such as identity-based encryption (IBE). Furthermore, to overcome the above security threats, such kind of identity-based access control placed on the shared data should meet the following security goals: • Data confidentiality: Unauthorized users should be prevented from accessing the plaintext of the shared data stored in the cloud server. In addition, the cloud server, which is supposed to be honest but curious, should also be deterred from knowing plaintext of the shared data. • Backward secrecy: Backward secrecy means that, when a user's authorization is expired, or a user's secret key is compromised, he/she should be prevented from accessing the plaintext of the subsequently shared data that are still encrypted under his/her identity. • Forward secrecy: Forward secrecy



## International Journal of Advanced Research in Science, Engineering and Technology

Vol. 6, Special Issue, August 2019

International Conference on Recent Advances in Science, Engineering, Technology and Management at Sree Vahini Institute of Science and Technology-Tiruvuru, Krishna Dist, A.P.

means that, when a user's authority is expired, or a user's secret key is compromised, he/she should be prevented from accessing the plaintext of the shared data that can be previously accessed by him/her.

**MOTIVATION** It seems that the concept of revocable identity-based encryption (RIBE) might be a promising approach that fulfills the aforementioned security requirements for data sharing. RIBE features a mechanism that enables a sender to append the current time period to the ciphertext such that the receiver can decrypt the ciphertext only under the condition that he/she is not revoked at that time period. As indicated in Figure 1, a RIBE-based data sharing system works as follows:

**Step 1:** The data provider (e.g., David) first decides the users (e.g., Alice and Bob) who can share the data. Then, David encrypts the data under the identities Alice and Bob, and uploads the ciphertext of the shared data to the cloud server.

**Step 2:** When either Alice or Bob wants to get the shared data, she or he can download and decrypt the corresponding ciphertext. However, for an unauthorized user and the cloud server, the plaintext of the shared data is not available.

**Step 3:** In some cases, e.g., Alice's authorization gets expired, David can download the ciphertext of the shared data, and then decrypt-then-re-encrypt the shared data such that Alice is prevented from accessing the plaintext of the shared data, and then upload the re-encrypted data to the cloud server again.



Fig. 1. A natural RIBE-based data sharing system

## II. RELATED WORK

## Revocable identity-based encryption

The concept of identity-based encryption was introduced by Shamir [13], and conveniently instantiated by Boneh and Franklin [14]. IBE eliminates the need for providing a public key infrastructure (PKI). Regardless of the setting of IBE or PKI, there must be an approach to revoke users from the system when necessary, e.g., the authority of some user is expired or the secret key of some user is disclosed. In the traditional PKI setting, the problem of revocation has been



# International Journal of Advanced Research in Science, Engineering and Technology

Vol. 6, Special Issue, August 2019

International Conference on Recent Advances in Science, Engineering, Technology and Management at Sree Vahini Institute of Science and Technology-Tiruvuru, Krishna Dist, A.P.

well studied [15], [16], [17], [18], [19], and several techniques are widely approved, such as certificate revocation list or appending validity periods to certificates. However,

there are only a few studies on revocation in the setting of IBE. Boneh and Franklin [14] first proposed a natural revocation way for IBE. They appended the current time period to the ciphertext, and non-revoked users periodically received private keys for each time period from the key authority. Unfortunately, such a solution is not scalable, since it requires the key authority to perform linear work in the number of non-revoked users. In addition, a secure channel is essential for the key authority and non-revoked users to transmit new keys. To conquer this problem, Boldyreva, Goyal and Kumar [20] introduced a novel approach to achieve efficient revocation. They used a binary tree to manage identity such that their RIBE scheme reduces the complexity of key revocation to logarithmic (instead of linear) in the maximum number of system users. However, this scheme only achieves selective security. Subsequently, by using the aforementioned revocation technique, Libert and Vergnaud [21] proposed an adaptively secure RIBE scheme based on a variant of Water's IBE scheme [22], Chen et al. [23] constructed a RIBE scheme from lattices. Recently, Seo and Emura [24] proposed an efficient RIBE scheme resistant to a realistic threat called decryption key exposure, which means that the disclosure of decryption key for current time period has no effect on the security of decryption keys for other time periods. Inspired by the above work and [25], Liang et al. [26] introduced a cloud-based revocable identitybased proxy re-encryption that supports user revocation and ciphertext update. To reduce the complexity of revocation, they utilized a broadcast encryption scheme [27] to encrypt the ciphertext of the update key, which is independent of users, such that only non-revoked users can decrypt the update key. However, this kind of revocation method cannot resist the collusion of revoked users and malicious non-revoked users as malicious nonrevoked users can share the update key with those revoked users. Furthermore, to update the ciphertext, the key authority in their scheme needs to maintain a table for each user to produce the re-encryption key for each time period, which significantly increases the key authority's workload

## III. IMPLEMENTATION

## **KUNodes algorithm**

Our RS-IBE scheme uses the same binary tree structure introduced by Boldyreva, Goyal and Kumar [20] to achieve efficient revocation. To describe the revocation mechanism, we first present several notations. Denote by  $\epsilon$  the root node of the binary tree BT , and Path( $\eta$ ) the set of nodes on the path from  $\epsilon$  to the leaf node  $\eta$  (including  $\epsilon$  and  $\eta$ ). For a non-leaf node  $\theta$ , we let  $\theta_l$  and  $\theta_r$  stand for its left and right

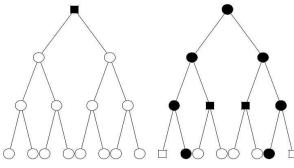


Fig. 2. An instance of the algorithm KUNodes

Informally, to identify the set Y, the algorithm first marks all the ancestors of revoked nodes as revoked, then outputs all the non-revoked children of revoked nodes. As an exam-ple, we present two instances of the algorithm KUNodes in Figure 2. The formal description is given below.

Algorithm 1 KUNodes(BT, RL, t)

1: X,Y←-Ø

2: for all  $(\eta_i, t_i) \in RL$  do

3: **if**  $t_i \le t$  **then** 

4: Add Path $(\eta_i)$  to X

5: **end if** 



## International Journal of Advanced Research in Science, Engineering and Technology

Vol. 6, Special Issue, August 2019

International Conference on Recent Advances in Science, Engineering, Technology and Management at Sree Vahini Institute of Science and Technology-Tiruvuru, Krishna Dist, A.P

6: end for

7: for all  $\theta \in X$  do

8: **if**  $\theta_1 \in X$  then

**9:** Add  $\theta_1$  to Y

10: end if

11: **if**  $\theta_r \in X$  then

**12:** Add  $\theta_r$  to Y

13: **end if** 

14: end for

15: if  $Y = \emptyset$  then

16: Add the root node  $\varepsilon$  to Y

17: end if

18: return Y

## IV. CONCLUSION

Cloud computing brings great convenience for people. Particularly, it perfectly matches the increased need of sharing data over the Internet. In this paper, to build a cost-effective and secure data sharing system in cloud computing, we proposed a notion called RS-IBE, which supports identity revocation and ciphertext update simultaneously such that a revoked user is prevented from accessing previously shared data, as well as subsequently shared data. Furthermore, a concrete construction of RS-IBE is presented. The proposed RS-IBE scheme is proved adaptive-secure in the standard model, under the decisional  $\ell$ -DBHE assumption. The comparison results demonstrate that our scheme has advantages in terms of efficiency and functionality, and thus is more feasible for practical applications.

## REFERENCES

- {1] L. M. Vaquero, L. Rodero-Merino, J. Caceres, and M. Lindner, "A break in the clouds: towards a cloud definition," ACM SIGCOMM Computer Communication Review, vol. 39, no. 1, pp. 50–55, 2008.
- [2] iCloud. (2014) Apple storage service. [Online]. Available: https://www.icloud.com/
- [3] Azure. (2014) Azure storage service. [Online]. Available: http://www.windowsazure.com/
- [4] Amazon. (2014) Amazon simple storage service (amazon s3). [Online]. Available: http://aws.amazon.com/s3/
- [5] K. Chard, K. Bubendorfer, S. Caton, and O. F. Rana, "Social cloud computing: A vision for socially motivated resource sharing," Services Computing, IEEE Transactions on, vol. 5, no. 4, pp. 551–563, 2012.
- [6] C. Wang, S. S. Chow, Q. Wang, K. Ren, and W. Lou, "Privacypreserving public auditing for secure cloud storage," Computers, IEEE Transactions on, vol. 62, no. 2, pp. 362–375, 2013.
- [7] G. Anthes, "Security in the cloud," Communications of the ACM, vol. 53, no. 11, pp. 16-18, 2010.
- [8] K. Yang and X. Jia, "An efficient and secure dynamic auditing protocol for data storage in cloud computing," Parallel and Distributed Systems, IEEE Transactions on, vol. 24, no. 9, pp. 1717–1726, 2013.
- [9] B. Wang, B. Li, and H. Li, "Public auditing for shared data with efficient user revocation in the cloud," in INFOCOM, 2013 Proceedings IEEE. IEEE, 2013, pp. 2904–2912.
- [10] S. Ruj, M. Stojmenovic, and A. Nayak, "Decentralized access control with anonymous authentication of data stored in clouds," Parallel and Distributed Systems, IEEE Transactions on, vol. 25, no. 2, pp. 384–394, 2014.